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SOLVING SYSTEMS OF EQUATIONS

WITH

APPLICATIONS TO OPTIMIZATION

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CHAPTER I

AN INTRODUCTION AND OVERVIEW

Optimization is a topic of intense current interest to mathematicians and computer scientists with broad potential for application in government and industry. The importance of maximum efficiency at minimum cost concerns all levels of operations and is just one of the many uses of optimization techniques. The translating of information into mathematical functions by applied mathematicians, with the assistance of modern computers, has made the analysis of various situations a much more manageable and useful instrument for the prosperity and enlightenment of our generation.

Nevertheless, the preponderance of research conducted on such a useful implement as optimization is impossible to cover within one prospectus. Therefore, the scope of this research has been focused on the investigation of iterative methods when applied to the solution of systems of equations. In general, all optimization techniques rely on iterations, for which convergence theory is incomplete and quite difficult to apply in multivariable cases.

The problem of optimization entails the analysis of a differentiable function of several variables. It is necessary to determine the values of the n unconstrained variables x_1, x_2, \dots, x_n which make the function f a maximum or a minimum. For a given function $f(x_1, x_2, \dots, x_n)$, the maxima and minima can only occur where the n equations $\frac{\partial f}{\partial x_i} = 0, i=1, 2, \dots, n$ are simultaneously satisfied, or at a boundary point of the domain of the variables. [5]

Thus, it does seem reasonable to consider the solution of systems of

equations. As a particular example, one can consider the following function:

$$f(x,y) = xy - \frac{x^3}{3} - \frac{y^3}{3} + 10. \quad (1)$$

By taking partial derivatives, the following system is obtained:

$$f_x(x,y) = y - x^2 = 0, \quad f_y(x,y) = x - y^2 = 0. \quad (2)$$

Simple algebra will secure two solutions to the system, (0,0) and (1,1).

It can be shown by analytical techniques that $f(1,1)$ is a maximum and that (0,0) does not produce an extremum. [6]

Sometimes it is useful to transform a system such as (2) into an equivalent fixed point problem. For example the point (1,1) is a fixed point for the function g , defined by the rule:

$$g(x,y) = (u,v) \\ u = x + \frac{x-2yx^2+y^2}{4xy-1} \quad v = y + \frac{y-2xy^2+x^2}{4xy-1}. \quad (3)$$

Thus an iterative scheme for approximating the solution might be devised by the use of g . One starts from a suitable initial point (x_0, y_0) and lets $(x_{n+1}, y_{n+1}) = g(x_n, y_n)$. In general the choice of an initial point is quite important. Ideally, the point should be close to the actual solution. The source of the choice of g as well as the motivation leading to its definition, will be presented in Chapter 3.

Applying the fixed point method above beginning with the initial

guess $x_0=y_0=10$, the following results are obtained, rounded to four digits.
(For actual calculations see Appendix I.)

n	x_n	y_n
1	5.263	5.263
2	2.908	2.908
3	1.756	1.756
4	1.227	1.227
5	1.036	1.036
6	1.001	1.001
7	1.000	1.000

Table 1

It is possible to explore some methods analogous to the above in considerable detail. In order for the author to construct a method, the theory for a restricted case, systems of linear equations, was consulted. Using geometric intuition, each equation was conceptualized as a plane, then the necessary procedure to compute the unique orthogonal projection from a point not on the plane to a point in the plane was determined. Subsequently, the method was rediscovered in Pizer's book. [8] Interestingly, the method is always convergent. (A proof can be found in the above source.)

This technique will be discussed in Chapter 2, and a generalization of it to certain non-linear systems is discussed in Chapter 3. Opportunities for further development and study of this method will also be discussed. Alternate approaches to the problem are treated in [4] and [7].

CHAPTER II

THE LINEAR CASE

Suppose that a system of equations in n variables is defined by n functions F_i , $i = 1, 2, \dots, n$. $F_i(x_1, \dots, x_n) = 0$. The system is linear if each F_i is a function of the form:

$$F_i(x_1, \dots, x_n) = a_{i1}x_1 + \dots + a_{in}x_n - b_i.$$

Alternatively, the system may be viewed as a matrix equation, $AX - B = 0$, and the vector $X = (x_1, \dots, x_n)$ may be regarded as a solution of the system.

The typical approaches to the linear case entail both algebraic and iterative techniques. Systems such as $AX = b$, with coefficients that are mostly nonzero, are called dense, and primarily are solved by algorithms based on Gaussian elimination. When A is square, sparse (most coefficients are zero) and of large order, iterative methods are the preferred method of solution. [1]

A thorough discussion justifies an investigation of Gaussian elimination, the Jacobian method and the Gauss-Seidel method. The chapter concludes with a discussion of an iterative technique developed by the author and her advisor, which resulted after concentration primarily on the iterative approaches. They were chosen in expectation that they might suggest generalization to non-linear problems.

It is important first to point out that, in general, methods for nonlinear simultaneous equations will extend to linear simultaneous equations, but the converse is not true. Therefore, the following

methods should be understood only to apply to the linear case.

Gaussian Elimination

Some readers will recall that Gaussian elimination involves the simplification of a system of equations using row reductions. These, applied to the augmented co-efficient matrix, result in reduction to a lower triangular matrix. The resulting system is then much easier to solve.

An example is the linear system

$$x_1 + 2x_2 + 4x_3 = 3$$

$$x_1 + \quad \quad 2x_3 = 0$$

$$2x_1 + 4x_2 + x_3 = 3$$

represented by the matrix

$$\left[\begin{array}{ccc|c} 1 & 2 & 4 & 3 \\ 1 & 0 & 2 & 0 \\ 2 & 4 & 0 & 3 \end{array} \right], \text{ which reduces to } \left[\begin{array}{ccc|c} 1 & 0 & 0 & -6/7 \\ 0 & 1 & 0 & 15/14 \\ 0 & 0 & 1 & 3/7 \end{array} \right],$$

The latter matrix represents a system equivalent to the original.

It is certainly trivial to identify the vector $\begin{cases} x_1 = -6/7 \\ x_2 = 15/14 \\ x_3 = 3/7 \end{cases}$ as a

solution to the system. [2] The method also works successfully on a system

with infinite solutions or with no solutions.

It should be noted that Gaussian elimination can become less practical for problems on a sufficiently large scale. When n is small, the method is unquestionably preferred. If $n = 10$, Gaussian elimination requires only 430 operations. Compare this to the number of operations required by Cramer's rule, another popular technique. For the case $n = 10$, the number of operations is 39,916,800. [1]

The Jacobian Method

The Jacobian method uses simultaneous displacements. At the end of each iteration, the components of an approximate solution x^i are replaced simultaneously with the corresponding components of a vector x^{i+1} . In general, for the equation $AX=b$

$$x_j^{i+1} = \frac{b_j}{A_{jj}} - \sum_{\substack{k=1 \\ k \neq j}}^n \frac{A_{jk}}{A_{jj}} x_k^i \text{ for } j = 1, 2, \dots, n;$$

where the terms A_{jk} are entries of A ; the terms b_j are components of B . That is, the first equation solves for x_1 , the second equation for x_2 , and so on, up through the n^{th} equation for x_n . The j^{th} equation produces a value for x_j^{i+1} by using the values of x_k^i on the right side of the equation.

The convergence of the method can sometimes be assured by first

making the matrix either row or column diagonally dominant. A matrix

A is said to be row diagonally dominant if

$$|a_{ii}| > \sum_{j \neq i}^n |a_{ij}| \quad \text{for all } i=1,2,\dots,n,$$

and column diagonally dominant if

$$|a_{ii}| > \sum_{j \neq i}^n |a_{ji}| \quad \text{for all } i=1,2,\dots,n. \quad [9]$$

It is important to point out that convergence is possible even if these conditions do not apply.

Consider,

$$x_1 + 6x_2 + 2x_3 = 15$$

$$x_1 + x_2 - 6x_3 = -3$$

$$6x_1 + x_2 + x_3 = 9.$$

In accordance with the above, the equations are rearranged as follows:

$$6x_1 + x_2 + x_3 = 9$$

$$x_1 + 6x_2 + 2x_3 = 15$$

$$x_1 + x_2 - 6x_3 = -3$$

Solve equation 1 for x_1

$$x_1 = 3/2 - (1/6)x_2 - (1/6)x_3$$

$$x_2 = 5/2 - (1/6)x_1 - (1/3)x_3$$

$$x_3 = 1/2 + (1/6)x_1 + (1/6)x_2.$$

Let $x^0 = (0,0,0)$

$$x^1 = \begin{bmatrix} 3/2 \\ 5/2 \\ 1/2 \end{bmatrix} \quad x^2 = \begin{bmatrix} 1 \\ 25/12 \\ 7/6 \end{bmatrix} \quad x^3 = \begin{bmatrix} 23/24 \\ 35/18 \\ 73/72 \end{bmatrix}$$

The sequence $\{x^i\}$ is seemingly converging. [8]

Gauss - Seidel Iteration

The Gauss - Seidel method applies essentially the same idea as the Jacobian. A key difference is that the computed values are employed as soon as possible to obtain each component x_j^i . In general, for $AX = B$ the formal equation for a stage of the Gauss - Seidel method is as follows:

$$x_j^{i+1} = \frac{b_j}{A_{jj}} - \sum_{k=1}^{j-1} \frac{A_{jk}}{A_{jj}} x_k^{i+1} - \sum_{k=j+1}^n \frac{A_{jk}}{A_{jj}} x_k^i$$

For example, application of this method to the above using initial vector $(0,0,0)$ yields:

$$\begin{aligned} x_1^1 &= 3/2 \\ x_2^1 &= 5/2 - 1/6 (3/2) - 1/3 (0) = 9/4 \\ x_3^1 &= 1/2 + 1/6 (3/2) + 1/6 (9/4) = 9/8. \end{aligned}$$

Thus, $x^1 = (3/2, 9/4, 9/8)$. Returning to equation one, this estimate can be used to compute x_1^2, x_2^2, x_3^2 , and so forth. [8]

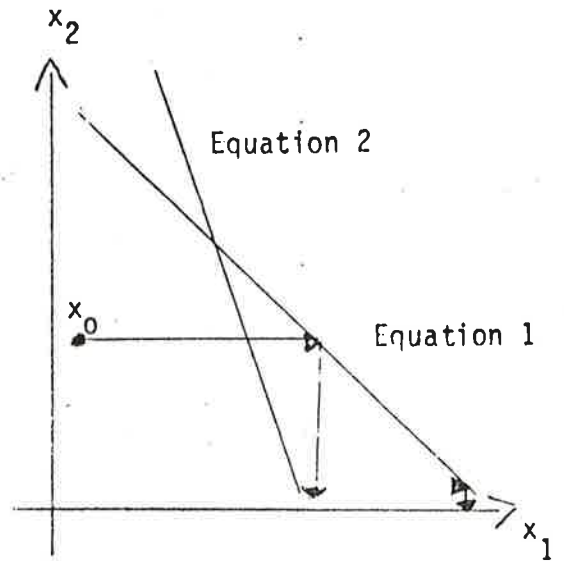
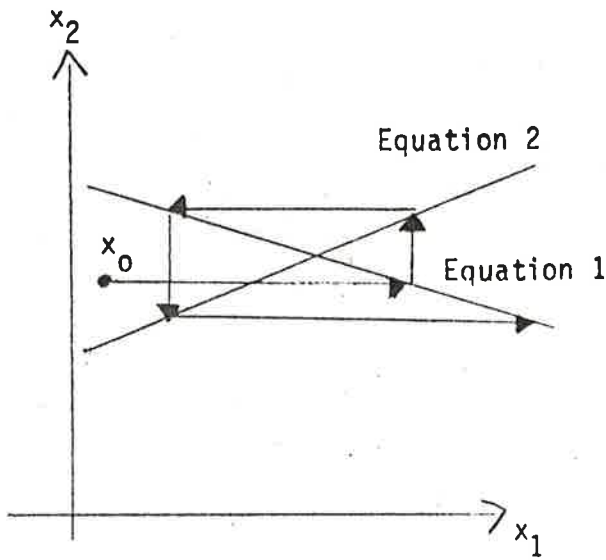
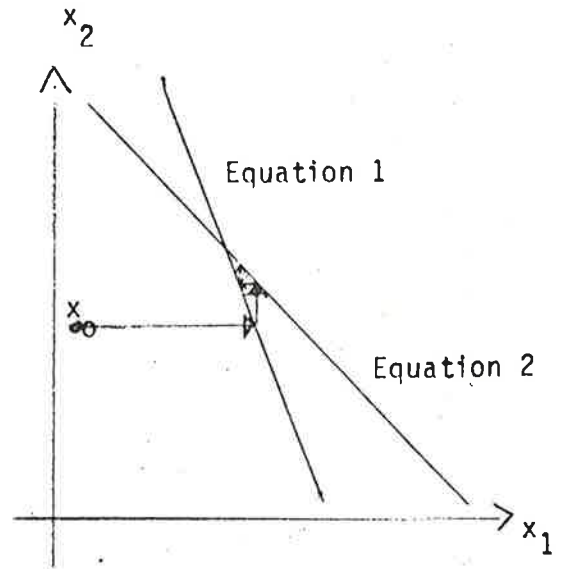
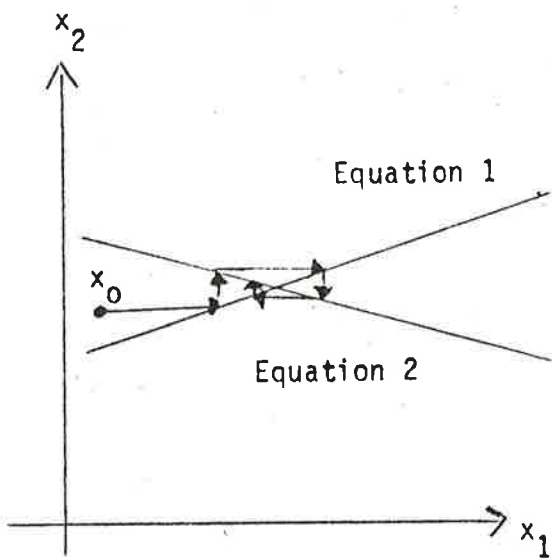


Figure 1

It might be noted at this point that convergence is guaranteed if both of the diagonal dominance conditions mentioned previously apply. Geometrically the convergence looks like the above. [8]

The Orthogonal Projection Method

The following method was developed by the author and her advisor. After initial development of the method, a thorough search of relevant literature resulted in the discovery that the method had previously been employed. A proof for convergence is found in Pizer's book. [8]

For clarity the discussion is confined to the case $n = 3$. The method extends easily to an arbitrarily large system.

In general, given a plane

$$a_1x_1 + a_2x_2 + a_3x_3 = b, \quad (1)$$

and a point $P = (u_1, u_2, u_3)$ not on the plane, the problem is to find the unique point $Q = (u'_1, u'_2, u'_3)$ on the plane such that the vector $P-Q$ is perpendicular to the plane.

$$P-Q = \langle u_1 - u'_1, u_2 - u'_2, u_3 - u'_3 \rangle \quad (2)$$

Since $P-Q$ is parallel to $\langle a_1, a_2, a_3 \rangle$, $\langle u_1 - u'_1, u_2 - u'_2, u_3 - u'_3 \rangle$ is equal to $C \langle a_1, a_2, a_3 \rangle$ for some constant scalar C .

$$\begin{array}{ll} u_1 - u'_1 = Ca_1 & u'_1 = u_1 - Ca_1 \\ u_2 - u'_2 = Ca_2 & \text{imply } u'_2 = u_2 - Ca_2 \\ u_3 - u'_3 = Ca_3 & u'_3 = u_3 - Ca_3. \end{array} \quad (3)$$

Replacing the variables in equation (1) by the coordinates of Q yields:

$$a_1 u_1' + a_2 u_2' + a_3 u_3' = b,$$

and combining (1) and (3)

$$a_1(u_1 - ca_1) + a_2(u_2 - ca_2) + a_3(u_3 - ca_3) = b. \quad (4)$$

Next solve for C.

$$\frac{a_1 u_1 + a_2 u_2 + a_3 u_3 - b}{a_1^2 + a_2^2 + a_3^2} = C. \quad (5)$$

Thus, the vector $\langle u_1', u_2', u_3' \rangle$ is obtained by using (5) in (3).

A FORTRAN Program (Appendix II) was used to implement the algorithm for various choices of n, giving the results in Table 2 for the system,

$$\begin{aligned} x_1 - x_2 + 2x_3 &= -3, \\ 5x_1 + \quad \quad 3x_3 &= 7, \\ 4x_1 - 3x_2 - x_3 &= 0. \end{aligned}$$

For brevity only one iterate from each cycle is included. The initial guess was (10, 10, 10). The iterates are rounded off to four places.

The exact solution is (2, 3, -1).

The method described above was apparently discovered by Dr. Lester

n	x_1	x_2	x_3
0	10.	10.	10.
3	6.9529	9.843	-1.7174
6	5.0673	7.9523	-3.5876
9	3.8999	6.2708	-3.2129
12	3.1768	5.0839	-2.5443
15	2.729	4.3073	-2.006
18	2.4516	3.8145	-1.6372
21	2.2797	3.5059	-1.3988
24	2.1733	3.3138	-1.2482
27	2.1073	3.1945	-1.154
30	2.0665	3.1205	-1.0955

Table 2

Levy of Long Island Jewish Medical Center. [8] Based on the Pythagorean theorem, the iteration moves closer and closer to the intersection of the planes. As mentioned earlier, the orthogonal projection method can be shown to be always convergent, provided a solution exists and is unique.

CHAPTER III

THE NON-LINEAR CASE

The Newton-Raphson method seems to be the proper point of departure into the nonlinear investigation. Based on a Taylor Series approach, the method for one variable expands to a multiple variable application.

Newton-Raphson

If f is continuous and twice differentiable, then

$$f(x_{i+1}) = f(x_i) + f'(x_i)(x_{i+1} - x_i) + \frac{f''(x_i)(x_{i+1} - x_i)^2}{2!} + \dots$$

If $f(x_{i+1})$ is close to zero, and x_i is very close to x_{i+1} so that $(x_{i+1} - x_i)^2$ and all higher powers can be neglected. then

$$0 \approx f(x_i) + f'(x_i)(x_{i+1} - x_i).$$

This suggests the equation:

$$x_{i+1} = x_i - \frac{f(x_i)}{f'(x_i)},$$

which is commonly called the Newton-Raphson method.

Given two continuous and twice differentiable functions, f_1 and f_2 , definition of two simultaneous equations is possible:

$$f_1(x,y) = 0; \quad f_2(x,y) = 0. \quad (1).$$

To generate a sequence $\{(x_i, y_i)\}$, which converges to a solution, one should use the two-dimensional Taylor series, expressing $f_1(x_{i+1}, y_{i+1})$ as

$$f_1(x_i, y_i) + \frac{\partial f_1}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_1}{\partial y} (y_{i+1} - y_i) + \dots$$

where the partial derivatives $\frac{\partial f_1}{\partial x}$ and $\frac{\partial f_1}{\partial y}$ are evaluated at $x = x_i$

and $y = y_i$. In the same way, f_2 is expanded as follows:

$$f_2(x_{i+1}, y_{i+1}) = f_2(x_i, y_i) + \frac{\partial f_2}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_2}{\partial y} (y_{i+1} - y_i) + \dots$$

If (x_{i+1}, y_{i+1}) and (x_i, y_i) are very close to the root (\bar{x}, \bar{y}) , the equations are simplified to:

$$0 \approx f_1(x_i, y_i) + \frac{\partial f_1}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_1}{\partial y} (y_{i+1} - y_i)$$

$$0 \approx f_2(x_i, y_i) + \frac{\partial f_2}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_2}{\partial y} (y_{i+1} - y_i).$$

Rearranging,

$$\frac{\partial f_1}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_1}{\partial y} (y_{i+1} - y_i) \approx -f_1(x_i, y_i) \quad (2)$$

$$\frac{\partial f_2}{\partial x} (x_{i+1} - x_i) + \frac{\partial f_2}{\partial y} (y_{i+1} - y_i) \approx -f_2(x_i, y_i).$$

For convenience, let $h = (x_{i+1} - x_i)$ and $k = (y_{i+1} - y_i)$. These equations can now be solved for h and k by the use of determinants, provided the determinant of the coefficient matrix does not vanish at any time during the iteration process. Therefore,

$$h = \frac{\begin{vmatrix} -f_1(x_i, y_i) & \frac{\partial f_1}{\partial y} \\ -f_2(x_i, y_i) & \frac{\partial f_2}{\partial y} \end{vmatrix}}{\begin{vmatrix} \frac{\partial f_1}{\partial x} & \frac{\partial f_1}{\partial y} \\ \frac{\partial f_2}{\partial x} & \frac{\partial f_2}{\partial y} \end{vmatrix}}, \quad k = \frac{\begin{vmatrix} \frac{\partial f_1}{\partial x} & -f_1(x_i, y_i) \\ \frac{\partial f_2}{\partial x} & -f_2(x_i, y_i) \end{vmatrix}}{\begin{vmatrix} \frac{\partial f_1}{\partial x} & \frac{\partial f_1}{\partial y} \\ \frac{\partial f_2}{\partial x} & \frac{\partial f_2}{\partial y} \end{vmatrix}}$$

The denominator is commonly called the Jacobian of the system.

After h and k are determined, x_{i+1} and y_{i+1} must be evaluated. To compute the next approximation set $x_i = x_{i+1}$, $y_i = y_{i+1}$ and repeat. [9]

The discerning reader might notice that the system (2) can be expressed in matrix notation.

$$J(x^I)(x^{I+1} - x^I) = -F(x^I), \quad (3)$$

where J is the Jacobian matrix for the vector valued function F , x^I is the vector (x_i, y_i) , and x^{I+1} is the next iterate (x_{i+1}, y_{i+1}) . If equation (3) is solved for x^{I+1} , the following is obtained:

$$x^{I+1} = x^I - J^{-1}(x^I) F(x^I). \quad (4)$$

The latter equation, which generalizes clearly to systems with an arbitrary number of variables, is often taken as a definition of Newton's method for iteratively solving systems of equations. Since the matrix J is difficult to invert, the usual procedure is to solve equation (3) for the vector $x^{I+1} - x^I$, as indicated earlier. The necessity of solving a matrix equation is one of the chief drawbacks of the method, especially as the number of unknowns increases.

Equation (4) also suggests the definition of a fixed point problem $G(x) = x$, which is equivalent to the problem $F(x) = 0$:

$$G(x) = x + J^{-1}(x) F(x). \quad (5)$$

For the example discussed in Chapter 1, the function F is defined by two equations:

$$f_1(x,y) = y - x^2, \quad f_2(x,y) = x - y^2.$$

The matrixes $J(x,y)$ and $J^{-1}(x,y)$ are given by:

$$J(x,y) = \begin{bmatrix} -2x & 1 \\ 1 & -2y \end{bmatrix}; \quad J^{-1}(x,y) = \begin{bmatrix} 2y/(1-4xy) & 1/(1-4xy) \\ 1/(1-4xy) & 2x/(1-4xy) \end{bmatrix}$$

By performing the computations indicated in (5) one obtains the following representation of G used in Chapter 1 to solve $F(x,y) = (0,0)$.

$$g_1(x,y) = x - \frac{x - 2yx^2 + y^2}{1-4xy}, \quad g_2(x,y) = y - \frac{y - 2xy^2 + x^2}{1-4xy}. \quad (7)$$

The Newton-Raphson method in practice is quite complicated. The requirement that the Jacobian is not allowed to be zero is difficult to predict without calculations and includes the risk that the process will be inoperative. In general, for n simultaneous equations in n unknowns, one must find n^2 partial derivatives, and must then solve for n increments h, k, \dots by solving n simultaneous equations for n unknowns. Therefore, mathematicians have sought an easier method, which, while perhaps not converging as quickly, is easier to apply. [9]

One modification of the Newton-Raphson method consists of applying the single variable Newton-Raphson method n times, once for each variable,

assuming the other variables are fixed.

In general, for two equations and two unknowns:

$$f_1(x,y) = 0 \quad \text{and} \quad f_2(x,y) = 0, \quad \text{and}$$

using x_0 and y_0 as initial approximations, one calculates a new value x_1 , from the single-variable Newton-Raphson method:

$$x_1 = x_0 - \frac{f_1(x_0, y_0)}{\frac{\partial f_1}{\partial x}}$$

where $\partial f_1 / \partial x$ is evaluated at x_0 and y_0 . Using f_2 and the most recent values of x and y , in this case x_1 and y_0 , a new value y_1 is calculated as follows:

$$y_1 = y_0 - \frac{f_2(x_1, y_0)}{\frac{\partial f_2}{\partial y}}$$

After obtaining x_1 and y_1 , one uses f_1 to calculate x_2 , and so on.

Note that the choice of using f_1 to calculate a new x and using f_2 to calculate a new y is not arbitrary. In fact, an improper choice of functions can lead to divergence. The function which has the steepest descent at the solution point should be used to find the next

x. [9]

There exists many other iterative methods designed to deal with single-variable, non-linear equations; the bisection method, the method of false position and various Picard iterations. However, methods like bisection and false position seem difficult to extend to many dimensions. The extension from single equations to systems of equations is a topic addressed briefly in most sources investigated by this researcher. One example, of course, is the modified Newton-Raphson method described above. Other techniques have been devised in various research papers. [4] and [10]

The following, an extension of the orthogonal vector method, is based partly on the underlying idea of the Gauss-Seidel iteration technique. The method apparently works with some effectiveness in several cases, though not in all of those attempted. Details of these results and potentially sufficient conditions for convergence will be discussed following the initial derivation.

Modified Orthogonal Vector Method

Suppose a system of equations is given,

$$F_i(x_1, \dots, x_n) = 0. \quad (8)$$

Suppose, in addition, that for each equation, there is a unique variable x_i , so that the particular equation can be solved for x_i in terms of the

other variables. Thus, one can write (8) in the form;

$$x_i = f_i(x_1, \dots, \hat{x}_i, \dots, x_n). \quad (9)$$

If $p_0 = (x_1^0, x_2^0, \dots, x_n^0)$, the initial point, and if q_1 is the point obtained in the surface F by varying only the first coordinate of p_0 , then:

$$\bar{x}_1 = f_1(x_2^0, \dots, x_n^0) \quad \text{and} \quad q_1 = (\bar{x}_1, x_2^0, \dots, x_n^0).$$

The plane tangent to F_1 at q_1 is represented analytically by:

$$x_1 - \bar{x}_1 = \frac{\partial f_1}{\partial x_2} (x_2 - x_2^0) + \dots + \frac{\partial f_1}{\partial x_n} (x_n - x_n^0). \quad (10)$$

In order to project orthogonally to a point p_1 in this plane, a normal vector is identified as:

$$n = \left(-1, \frac{\partial f_1}{\partial x_2}, \dots, \frac{\partial f_1}{\partial x_n} \right).$$

Then $p_1 - p_0 = cn$ (for some constant c).

More specifically, if $p_1 = (x_1^1, \dots, x_n^1)$ then:

$$x_1^1 - x_1^0 = -c, \quad \text{and} \quad x_i^1 - x_i^0 = c \frac{\partial f_1}{\partial x_i} \quad \text{for } 2 \leq i \leq n. \quad (11)$$

Isolating the coordinates of p_1 yields:

$$x_1^1 = x_1^0 - c; \quad x_i^1 + c \frac{df_1}{dx_i} \quad \text{for } 2 \leq i \leq n. \quad (12)$$

To solve for c , one uses the equation of the tangent plane (10) and the system (11):

$$x_1^1 - \bar{x}_1 = x_1^0 - c - \bar{x}_1 = \frac{\partial f_1}{\partial x_2} (x_2^1 - x_2^0) + \dots + \frac{\partial f_1}{\partial x_n} (x_n^1 - x_n^0).$$

$$\begin{aligned} x_1^0 - \bar{x}_1 &= c + \frac{\partial f_1}{\partial x_2} \left(c \frac{\partial f_1}{\partial x_2} \right) + \dots + \frac{\partial f_1}{\partial x_n} \left(c \frac{\partial f_1}{\partial x_n} \right) \\ &= c \left\{ 1 + \left(\frac{\partial f_1}{\partial x_2} \right)^2 + \dots + \left(\frac{\partial f_1}{\partial x_n} \right)^2 \right\} \end{aligned}$$

$$c = \frac{x_1^0 - \bar{x}_1}{1 + \left\{ \left(\frac{\partial f_1}{\partial x_2} \right)^2 + \dots + \left(\frac{\partial f_1}{\partial x_n} \right)^2 \right\}} \quad (13)$$

Now, the iterate p_1 is computed from (12).

To obtain a general iterative step; $1 \leq i \leq n$, one follows an entirely similar procedure.

$$\begin{aligned} \text{If } p_i &= (x_1^i, \dots, x_n^i) \\ \bar{x}_i^i &= f_i(x_1^i, \dots, x_i^i, \dots, x_n^i) \end{aligned}$$

$$c = \frac{x_i^i - x_i^{i-1}}{\quad} \quad \text{where}$$

$$1 + \left(\frac{\partial f}{\partial x_i}\right)^2 + \dots + \left(\frac{\partial f}{\partial x_{i-1}}\right)^2 + \left(\frac{\partial f}{\partial x_{i+1}}\right)^2 + \left(\frac{\partial f}{\partial x_n}\right)^2$$

all partials are evaluated at $(x_1^i, \dots, x_i^i, \dots, x_n^i)$.

$$x_i^{i+1} = x_i^i - c, \text{ and for } j \neq i, x_j^{i+1} = x_j^i + c \frac{\partial f}{\partial x_j}$$

If $i > n$, then $i = nq + r$ for some $r, 0 \leq r < n, q \geq 1$. Thus p_{i+1} can be generated from p_{nq} in the same way that p_i is generated from p_0 above when $i \leq n$.

The method was applied on several examples, producing the following results. The algorithm was terminated when $\sum_{i=qn+1}^{qn+n} \|p_i - p_{i-1}\| < \epsilon$.

The sum is taken over a cycle of n iterates. $\epsilon = .1$.

Example 1

This problem arose in Chapter 1 in connection with the problem of determining the critical points for the function:

$$f(x,y) = xy - \frac{x^3}{3} - \frac{y^3}{3} + 10.$$

$$\begin{cases} x = y^2 \\ y = x^2 \end{cases}$$

n	x_n	y_n	n	x_n	y_n
0	1.	2.	0	.6	.6
1	1.326	1.5318	1	.7348	.4856
2	1.1748	1.318	2	.731	.4704
3	1.1	1.1898	3	.7161	.4475
4	1.0598	1.1156	4	.6863	.4051
5	1.0365	1.0714	5	.6213	.3234
			6	.4575	.1691
			7	.1144	.0113
			8	1.928E-04	3.7126E-08
			9	0.0	0.0

Table 3

Example 2

$$\begin{cases} x = y - y^3 \\ y = x^3 - x \end{cases}$$

n	x_n	y_n
0	10.	10.
1	6.7039	6.6666
2	4.5248	4.4436
3	3.099	2.9595
4	2.1868	1.9823
5	1.6271	1.2709
6	1.2909	.7173
7	.5297	-.382
8	-.0532	.0544
9	1.5865E-04	-1.5866E-04
10	-4.02 E-12	4.02 E-12

Table 4

Example 2A

A modification of Example 2 to:

$$\begin{cases} x = 2(y-y^3) \\ y = 2(x^3-x), \end{cases}$$

does not converge from $p_0 = (10,10)$, but does converge from $p_0 = (.4,.4)$.

Example 3 This problem known as Rosenblock's parabolic valley is found in [4].

$$F(x,y) = 100(y-x^2)^2 + (1-x)^2$$

$$\begin{cases} x = y \\ y = x^2 - 1/200x + 1/200 \end{cases} \quad \epsilon = .001$$

n	x_n	y_n	n	x_n	y_n
0	1.2	1.0	0	0.86	1.14
1	1.0393	1.803	1	1.0297	1.0605
2	1.0393	1.803	2	1.0297	1.0605

Table 5

Example 4 [4]

$$F(x,y,z) = (x-y+z)^2 + (y+z-x)^2 + (x+y-z)^2$$

$$\begin{cases} x = (y+z)/3 \\ y = (x+z)/3 \\ z = (x+y)/3 \end{cases}$$

n	x _n	y _n	z _n	n	x _n	y _n	z _n
0	.5	1.	.5	0	100.	-1.	2.5
1	.4038	.6574	.554	1	26.5984	23.1687	16.589
2	.3352	.5631	.4424	2	23.2201	18.1613	13.7938
3	.2772	.4637	.3678	3	19.0713	15.1385	11.4033
4	.2293	.3838	.304	4	15.7928	12.5107	9.4345
5	.1897	.3175	.2516	5	13.0632	10.3512	7.8048
6	.1569	.2627	.2081	6	10.807	8.5631	6.4567
7	.1298	.2173	.1722	7	8.9403	7.084	5.3415
8	.1074	.1798	.1424	8	7.396	5.8604	4.4188
9	.0888	.1487	.1178	9	6.1186	4.8482	3.6556
10	.0735	.123	.0974	10	5.0617	4.0108	3.0242
11	.0608	.1018	.0806	11	4.1874	3.318	2.5018
12	.0503	.0842	.0667	12	3.4541	2.7449	2.0697
13	.0416	.0697	.0552	13	2.8658	2.2708	1.7122
14	.0344	.0576	.0456	14	2.3708	1.8786	1.4164
15	.0284	.0477	.0378	15	1.9613	1.5541	1.1718
16	.0236	.0394	.0312	16	1.6225	1.2856	.9694
17	.0348	.0275	.0207	17	1.3422	1.0636	.802
				18	1.11	.8799	.6634
				19	.9186	.7279	.5433
				20	.75	.6022	.454
				21	.6287	.4981	.3756
				22	.52	.4121	.3107
				23	.4304	.3409	.257
				24	.3559	.282	.2127
				25	.2945	.2333	.1759
				26	.2436	.1930	.1455
				27	.2015	.1597	.1204
				28	.1667	.1321	.0996
				29	.1379	.1093	.0824
				30	.1141	.0904	.0682
				31	.0944	.0748	.0564
				32	.0781	.0619	.0467
				33	.0646	.0512	.0386
				34	.0534	.0423	.0319
				35	.0442	.035	.0264
				36	.0366	.029	.0219

Table 6

Convergence and Efficiency Considerations

At this writing, attempts to obtain sufficient conditions for convergence of the preceding iteration technique have been unsuccessful. It is strongly suspected that convergence can be guaranteed by the choice of an initial point P_0 sufficiently close to the solution P , provided all of the first and second order partial derivatives of the functions f_i are continuous at P . The computational results above lend considerable support to this conjecture. Note, the method is a fixed point type of iteration. I.e., a function g so that $p_{n+1} = g(p_n)$ is obtainable for each iterate p_n . Thus, a well known convergence theorem might be useful in particular cases. The following theorem is taken from [8].

If g_i denotes the coordinate functions defining g , the following can be assumed.

Theorem Let K be a positive constant with $K < 1$. Let p be a fixed point for the function g , and let R be a positive constant. If

$$\sum_{k=1}^n \left| \frac{\partial g_i}{\partial x_k}(x) \right| \leq K \text{ for all } x \text{ satisfying } |p - x|_{\infty} \leq R. \text{ Then the}$$

iteration given by $p_{n+1} = g(p_n)$ converges to p for any choice of initial point satisfying $|p_0 - p|_{\infty} \leq R$.

For a slightly different version of this result see [3].

In order to make a rough comparison between the efficiency of the orthogonal projection method and Newton's method, the number of arithmetic operations required for the projection method was checked. To compute one

cycle of iterates, p^{mn} , p^{mn+1} , . . . , p^{mn+n} the number of operations required is $4n^2-n$. In addition $n^2 + n$ function evaluations are required. By comparison Newton's method requires the same number of function evaluations as well as the solution of a linear system of n equations in n variables. As n increases in size, the speed of machine operation might become a factor in favor of the projection algorithm, since the number of operations needed to solve the linear system of n equations is $O(n^3)$.

The computer implementation used by the author to test the projection method is indicated in Appendix III. The storage requirements for this program are roughly the same as would be needed for an implementation of Newton's method. Nevertheless, it now seems that the program is less efficient in this regard than it might be. For example, it does not seem necessary to require space for all of the n^2 partial derivatives at any one time. At most n cells should be sufficient for each iterative step. Another economy might be obtained by storing each of the function values $f_i(x_1^m, \dots, \hat{x}_i^m, \dots, x_n^m)$ in the same cell. These considerations would become significant if the algorithm is applied to large-scale problems such as those commonly encountered in the numerical solutions of partial differential equations. It is acknowledged, of course, that the number of function evaluations limits this possibility significantly.

Conclusion

In conclusion, it is necessary to mention certain aspects of this work which are believed to warrant further study. Due to time restraints, the algorithm was not applied to a sufficiently wide variety of problems. Thus, further testing of the algorithm would provide useful insight into matters of efficiency, speed and storage requirements. In addition, this information would provide adequate data for a comparison between the method and other well known techniques, such as Newton's method.

A potential limitation of the orthogonal method is the necessity to solve for x_j in each equation and compute derivatives for the explicit function obtained thereby. In some problems, neither of these steps is practicable. Perhaps an application of the one variable version of Newton's method for approximating x_j in conjunction with the use of implicit differentiation for evaluation of partial derivatives would extend applicability of the algorithm to some of these more difficult problems.

There are some aspects of optimization, currently of great interest, have not been touched upon at all. This is perhaps, best illustrated by the fact that these solutions or results do not in any case help to determine if an extremum is located at the point in question. Nevertheless, the functions to be optimized usually exhibit behavior near a critical point which can help to identify the point as an extremum. For example, the directional derivatives may change in sign at the point, or the second order partial derivatives may be positive definite near the point. Such features

might be useful in the construction of types of algorithms which are totally different from the approach of this paper. For further information, see [4] and [5].

Needless to say, the facts are not all here. The author has left many open questions which require further research. Hopefully, this work will be extended beyond its present stance, either by the author and/or an intrigued reader.

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APPENDIX

APPENDIX I

$$x + \frac{x-2yx^2+y^2}{4xy-1}$$

$$y + \frac{y-2xy^2+x^2}{4xy-1}$$

$$x_0=y_0=10$$

$$x_1=y_1=10 + \frac{10-2(10)^3+10^2}{4(10)^2-1}$$

$$=10 - 4.7368$$

$$=5.2632$$

$$x_2=y_2=5.2632 + \frac{5.2632-2(5.2632)^3+(5.2632)^2}{4(5.2632)^2-1}$$

$$=5.2632 - 2.355$$

$$=2.9082$$

$$x_3=y_3=2.9082 + \frac{2.9082-2(2.9082)^3+(2.9082)^2}{4(2.9082)^2-1}$$

$$=2.9082 - 1.1521$$

$$=1.7559$$

$$x_4=y_4=1.7559 - .5284$$

$$=1.2274$$

$$x_5=y_5=1.2274 - .19186$$

$$=1.0355$$

$$x_6=y_6=1.0355 - .0343$$

$$=1.001$$

$$x_7=y_7=1.001 - .000998$$

$$=1.000$$

APPENDIX II

```

C   THIS PROGRAM COMPUTES THE POINT OF INTERSECTION OF N PLANES IN R1,
    BY ORTHOGONAL PROJECTIONS INITIATED BY AN INITIAL GUESS

C   ESTABLISH SIZE OF ARRAYS
    DIMENSION P(11,10), A(10,10), B(10), D(10)

C   READ IN NUMBER OF PLANES BEING CONSIDERED
    READ (1,600) N

C   READ IN ERROR ALLOWANCE AND MAXIMUM NUMBER OF ALLOWED ITERATIONS
    READ (1,610) E, MAXIT
    M=0
    WRITE (3,620) N, E, MAXIT

C   READ IN AND WRITE INITIAL GUESS
    DO 10 J=1,N
    READ (1,630) P(1,J)
10  CONTINUE
    WRITE (3,640) (P(1,J),J=1,N)

C   READ AND WRITE EQUATIONS
    DO 20 I=1,N
    DO 30 J=1,N
    READ (1,650) A(I,J)
30  CONTINUE
    READ (1,660) B(I)
    WRITE (3,670) (A(I,J),J=1,N),B(I)
20  CONTINUE

C   COMPUTE DENOMINATOR
    DO 45 I=1,N
    D(I)=0.0
    DO 40 J=1,N
    G=A(I,J)
    D(I)=(D(I)+A(I,J)**2)
40  CONTINUE
45  CONTINUE

C   COMPUTE C
    I=1
100 F=0.0
    DO 50 J=1,N
    H=A(1,J)
    F=F+H*P(I,J)
50  CONTINUE
    C=(F-B(I))/D(I)

C   COMPUTE ORTHOGONAL PROJECTION POINT AND REPLACE INITIAL GUESS WITH IT
    DO 60 J=1,N
60  P(I+1,J)=P(I,J)-A(I,J)*C
    WRITE (3,680) (P(I+1,J),J=1,N)

```

APPENDIX II (Continued)

```

C   CHECK TO SEE IF NEW PROJECTION POINT HAS BEEN COMPUTED FOR EACH PLANE
    IF (I .EQ. N) GO TO 200
C   IF NOT RETURN FOR MORE DATA
    I=I+1
    GO TO 100

C   FIRST CHECK TO SEE IF MAXIMUM NUMBER OF ITERATIONS HAS OCCURRED
200 M=M+1
    IF (M .GE. MAXIT) STOP

C   CHECK TO SEE IF CONCLUSION REACHED
    S=0.0
    DO 300 I=1,N
    DO 300 J=1,N
    S=S+ ABS(P(I,J) - P(I+1,J))
300 CONTINUE
    IF (S .GE. E) THEN DO
    DO 400 J=1,N
    I=1
    P(I,J)=P(N+1,J)
400 CONTINUE
    GO TO 100
    ELSE DO
10  WRITE (3,700) (P(N+1,J),J=1,N)
    END IF

```

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APPENDIX III

```
40 DIM R(10): DIM P(10,10): DIM F(10): DIM D(10,10): DIM R1(10)
45 PRINT "SELECT ONE OF SUBROUTINES"
46 INPUT Q
50 PRINT "ENTER NUMBER OF EQUATIONS BEING CONSIDERED"
60 INPUT N
70 PRINT "ENTER ERROR ALLOWANCE "
80 INPUT E
90 PRINT "ENTER MAXIMUM NUMBER OF ITERATIONS ALLOWED"
100 INPUT M
110 PRINT
120 FOR J = 1 TO N
130 PRINT "ENTER INITIAL GUESS", J
140 INPUT R(J)
150 LET R1(J) = R(J)
160 NEXT J
170 PRINT

180 REM "SET UP COUNTER TO MONITOR NUMBER OF ITERATIONS"
185 LET K = 0
190 LET I = 0
200 I = I + 1
210 PRINT

220 REM "THE SUBROUTINE MUST HAVE BEEN CALCULATED AND ENTERED PREVIOUSLY
    FOR EACH NON-LINEAR EQUATION"
230 GOSUB 620
235 LET X = R(I)
240 LET Y = F(I)
250 PRINT

260 REM "COMPUTE SCALAR C"
270 LET S = 1
280 FOR J = 1 TO N
290 IF J = I THEN 300
292 LET S = S + D(I,J) **.2
300 NEXT J
310 LET C = (X - Y) / S
320 PRINT

330 REM "COMPUTE INDIVIDUAL CORDINATE OF PROJECTION POINT"
340 FOR J = 1 TO N
350 IF J = I THEN 355
352 LET P(I,J) = R(J) + C * D(I,J)
354 GO TO 360
355 LET P(I,J) = R(J) - C
360 NEXT J
370 FOR J = 1 TO N
380 PRINT P(I,J),
385 LET R(J) = P(I,J)
390 NEXT J
400 PRINT
```

APPENDIX III (Continued)

```

405 REM "CHECK IF NEW PROJECTION POINT HAS BEEN COMPUTED FOR EACH SURFACE"
410 IF I = N THEN 425
420 GO TO 200
425 LET K = K + 1

426 REM "IF MAXIMUM NUMBER OF ITERATIONS HAVE BEEN COMPUTED STOP"
430 IF K = M THEN 550

435 REM "CHECK TO SEE IF CONCLUSION REACHED"
437 LET T = 0.0
440 FOR J = 1 TO N
445 LET I = 1
447 LET T = T + ABS (R1(J) - P(1,J))
450 FOR I = 2 TO N
460 T = T + ABS (P(I,J) - P(I-1,J))
470 NEXT I
480 NEXT J
500 IF T = E THEN 560
510 PRINT "POINT OF INTERSECTION IS"
520 FOR J = 1 TO N
530 PRINT P(N,J)
540 NEXT J
543 PRINT
544 PRINT
545 PRINT "NUMBER OF CYCLES", K
550 STOP
560 FOR J = 1 TO N
570 LET R1(J) = R(J)
580 NEXT J
590 GO TO 190
620 IF Q = 1 THEN 1000
630 IF Q = 2 THEN 900
640 IF Q = 3 THEN 960
650 IF Q = 4 THEN 970

660 REM "NEXT SUBROUTINE GOES HERE"
900 GO SUB 2000
950 RETURN
960 GO SUB 3000
965 RETURN
970 GO SUB 4000
975 RETURN
1000 IF I = 1 THEN 1300
1100 LET F(I) = R(1) **2
1200 LET D(I,1) = 2 * R(1)
1205 RETURN
1300 LET F(I) = R(2) **2
1400 LET D(I,2) = 2 * R(2)
1405 RETURN

```

APPENDIX III (Continued)

```

2000 IF I = 1 THEN 2300
2100 LET F(I) = R(1) **3 - R(1)
2200 LET D(I,1) = 3 * R(1) **2 - 1
2250 RETURN

2300 LET F(I) = R(2) - R(2) **3
2400 LET D(I,2) = 1 - 3 * R(2) ** 2
2500 RETURN

3000 IF I = 1 THEN 3300
3100 LET F(I) = R(1) **2 - 1 (200 * R(1)) + 1 / 200
3200 LET D(I,1) = 2 * R(1) + 1 / (200 * R(1) **2)
3205 RETURN

3300 LET F(I) = SQR (R(2))
3400 LET D(I,2) = 1 / (2 * (SQR (R(2))))
3401 RETURN

4000 IF I = 1 THEN 4300
4001 IF I = 2 THEN 4500
4100 LET F(I) = (R(1) + R(2)) / 3
4200 LET D(I,1) = 1/3
4205 LET D(I,2) = 1/3
4210 RETURN

4300 LET F(I) = (R(2) + R(3))/ 3
4400 LET D(I,2) = 1/3
4405 LET D(I,3) = 1/3
4406 RETURN

4500 LET F(I) = R(1) + R(3)) / 3
4600 LET D(I,1) = 1/3
4700 LET D(I,3) = 1/3
4701 RETURN
4702 END

```